A Class of Sums with Unexpectedly High Cancellation

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June 3, 2020

Pentagonal Number Theorem

Let p(n) be the number of partitions of n and $G_\ell = \frac{\ell(3\ell-1)}{2}$ be $\ell-$ th pentagonal number. Then

$$\sum_{G_{\ell} < n} (-1)^{\ell} \rho(n - G_{\ell}) = 0.$$

Proof in for example Professor Berndt's "Number Theory in Spirit of Ramanujan" Book.

Rademacher expression for p(n)

Let $\mu_k(n) = \frac{\pi\sqrt{24n-1}}{6k}$. Rademacher-Ramanujan-Hardy proved that

$$p(n) = \frac{\sqrt{12}}{24x - 1} \left(\sum_{k=1}^{\infty} A_k(n) \left(\left(1 - \frac{1}{\mu_k(n)}\right) e^{\mu_k(n)} + \left(1 + \frac{1}{\mu_k(n)}\right) e^{-\mu_k(n)} \right) \right)$$

where

$$A_k(n) = \sum_{\substack{0 \le h < k \\ (h,k)=1}} \omega_{h,k} e^{\frac{2\pi i h n}{k}}.$$

Proof in for example Professor Andrew's "Theory of partitions" book. Approximate version

$$p(n) \simeq \frac{e^{\pi\sqrt{\frac{2n}{3}}}}{4\sqrt{3}n}.$$

Approximation of number of partitions

The first two terms:

$$p_2(x) = \frac{\sqrt{12}e^{\frac{\pi}{6}\sqrt{24x-1}}}{24x-1}\left(1 - \frac{6}{\pi(24x-1)^{\frac{1}{2}}}\right) + O(p(x)^{0.5}).$$

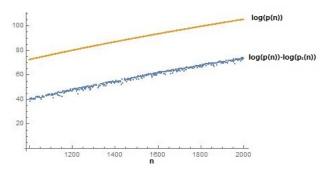


Figure: Comparison of the error term of first two terms with actual number for 20 < n < 2000.

Conclusion

What will happen if we use $p_2(n)$ in pentagonal number theorem?

$$\sum_{G_{\ell} \leq n} (-1)^{\ell} p_2(n - G_{\ell}) = O(p(n)^{0.5}).$$

Beginning of a long story!

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Observation

Lets test something simpler! We proved that

$$\sum_{l^2 < n} (-1)^l e^{\sqrt{n-l^2}} = O(e^{\frac{\sqrt{n}}{100}}).$$

There is something deeper than a combinatorial property.

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Observation

Lets test something simpler!

$$\sum_{\ell^2 < n} (-1)^{\ell} e^{\sqrt{n-\ell^2}} = O(n^{10}).$$

Our estimated error term is very small!

Theorem

Let $b, d \in \mathbb{R}$, a, c > 0; Also, let h(x) be $(\alpha x + \beta)^t$ for $\alpha, \beta, t \in \mathbb{R}$. Then

$$\sum_{n:an^2+bn+d
(1)$$

where w > 0 is defined as follows. Set

$$\Delta(r) := \sqrt{\sqrt{a}r\frac{\sqrt{ar^2 + 4} + r\sqrt{a}}{2} - \frac{\pi r}{c}} \quad , \quad r \ge 0$$

If $r = \alpha$ is when $\Delta(r)$ is maximized, then $w = \min(1, \Delta(\alpha))$.

A heuristic argument

Consider Bernoulli random variables $\epsilon_n=\pm 1$ with probability $P(\epsilon_n=1)=0.5$. Then what is expectation of

$$E\left(\sum_{\ell^2 < n} \epsilon_{\ell} e^{\sqrt{n - \ell^2}}\right) = 0$$

$$Var\left(\sum_{\ell^2 < n} \epsilon_{\ell} e^{\sqrt{n - \ell^2}}\right) \gg e^{2\sqrt{n}}.$$

Then why this sum is that small?

First Natural try

Understanding using Taylor expansions:

$$\sum_{r=0}^{\infty} \frac{S_r(M)}{r!} := \sum_{r=0}^{\infty} \sum_{\ell^2 < 4M^2} (-1)^{\ell} \frac{(4M^2 - \ell^2)^{\frac{r}{2}}}{r!} = O(e^{\frac{2M}{50}}).$$

We expect that $deg(S_r(M)) = 2r$.

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Reality

We proved that actually $\deg(S_r(M)) = r - 1$.

$$S_4(M) = 16M^3 - 17M$$

 $S_6(M) = -408M^5 - 480M^3 - 2073M$.

This might be doable by using the known result about Bernoulli numbers. We did not attempt to prove it this way; we predict that this way can prove this polynomial case in the best case scenario.

Remark

These sums are similar to Kloosterman's sum, I call them sisters! Our interested series:

$$\sum_{\ell^2 \le n} (-1)^\ell e^{\sqrt{n-\ell^2}}$$

Kloosterman's sum:

$$\sum_{\substack{0 \leq \ell < n \\ \gcd(n,\ell) = 1}} e^{2\pi i \ell^{-1} (a - \ell^2)}$$

We prove a weaker result about their father:

$$\sum_{\ell^2 < Tx} (-1)^{\ell} e^{(\alpha+i\beta)\sqrt{x-\frac{\ell^2}{T}}} = O\left(\sqrt{\frac{Tx}{|\beta|+1}} e^{\alpha(\sqrt{\frac{2}{2+\pi^2}}+\delta)\sqrt{x}} + \sqrt{T}\right).$$

Theorem

Assume $\epsilon>0$, x is large enough and $a=1-\sqrt{\frac{2}{2+\pi^2}}$. We have

$$\sum_{\ell^2 < xe^{\frac{4a}{3}\sqrt{x}}} (-1)^\ell \Psi\left(e^{\sqrt{x-\ell^2 e^{-\frac{2a}{3}\sqrt{x}}}}\right) = \textit{O}\left(e^{(1-\frac{a}{3}+\epsilon)\sqrt{x}}\right).$$

Remarks

In particular for $T = e^{0.786\sqrt{x}}$:

$$\sum_{0 < 2\ell < \sqrt{xT}} \Psi \big([e^{\sqrt{x - \frac{(2\ell)^2}{T}}}, \ e^{\sqrt{x - \frac{(2\ell - 1)^2}{T}}}] \big) \ = \Psi \big(e^{\sqrt{x}} \big) \left(\frac{1}{2} + \textit{O} \left(e^{-0.196\sqrt{x}} \right) \right),$$

This says we have half of primes in $\cup [e^{\sqrt{x-\frac{(2\ell)^2}{T}}}, e^{\sqrt{x-\frac{(2\ell-1)^2}{T}}}]$. If we use RH naively, we cannot make the error term this small. So controlling the father can be really rewarding!!!

General case

In general if we consider a "Meinardus type" partition as

$$\lambda(n) \sim (g(n))^q e^{(k(n))^{\theta}} \left(1 - \frac{1}{(h(n))^r}\right) + O(\lambda(n)^s)$$

where 0 < s < 1 and $\theta, r, q > 0$ and k(n) is a linear polynomial and g(n), h(n) are rational functions. Then

$$\sum_{t(\ell) < n} (-1)^{\ell} \lambda(n - t(\ell)) = O(\lambda^{\kappa}(n))$$

where $\kappa > s$ is determined from the analytic properties of the partition.

Application in Partitions

We proved that for the usual partitions

$$\sum_{\ell^2 < x} (-1)^{\ell} p(x - \ell^2) \sim 2^{-3/4} x^{-1/4} \sqrt{p(x)}.$$

and for the partitions with distinct parts

$$\sum_{\ell^2 < n} (-1)^{\ell} q(n - \ell^2) = O(\sqrt[3]{q(n)})$$

The problem asks for biggest possible k < n such that there exists two disjoint sets $\{y_i\}$ and $\{x_j\}$ such that

$$x_1 + x_2 + \dots + x_n = y_1 + y_2 + \dots + y_n$$

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We can suggest a solution to an "approximate" version of PTE problem.

Theorem

Let $n \ge 1$ and define $N = \lfloor (2n)^{\frac{2}{3}} \rfloor$. Let for $1 \le i \le n$

$$x_i = N^3 - (2i - 2)^2 \in \mathbb{N}$$
 $y_i = N^3 - (2i - 1)^2 \in \mathbb{N}$

Then for all $1 \le r \ll k := \frac{n^{\frac{2}{3}}}{\log(n)}$ we have

$$\sum_{1 \le i \le n} x_i^r - \sum_{1 \le i \le n} y_i^r = O\left(\left(\max(x_i, y_i) \right)^{\frac{5r}{6}} \right) = O(N^{\frac{5r}{2}}). \tag{2}$$

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So modulo error $O\left((\max(x_i,y_i))^{\frac{5r}{6}}\right)$ we have $k\gg n^{\frac{2}{3}+\epsilon}$.

Proof using circle method

We just talk about the case $\sum_{\ell^2 < n} (-1)^{\ell} e^{\sqrt{n-\ell^2}}$. Let

$$f(z) = \frac{e^{\sqrt{n-z^2}}}{\sin(\pi z)}$$

and find $\int_{\gamma} f(z)dz$.

By the Residue Theorem the residue will be exactly the sum we are interested.

$$\int_{\gamma} f(z)dz = \sum_{\ell^2 \le n} (-1)^{\ell} e^{\sqrt{n-\ell^2}}.$$

We need to compute the integral over contour.

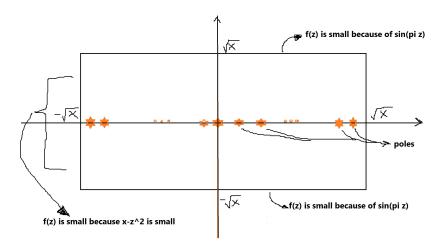


Figure: Contour γ

Thank You

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